

# FACEGroup: Feasible and Actionable Counterfactual Explanations for Group Fairness

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**Abstract.** Counterfactual explanations assess unfairness by revealing how inputs must change to achieve a desired outcome. This paper introduces the first graph-based framework for generating group counterfactual explanations to audit group fairness, a key aspect of trustworthy machine learning. Our framework, FACEGroup (Feasible and Actionable Counterfactual Explanations for Group Fairness), models real-world feasibility constraints, identifies subgroups with similar counterfactuals, and captures key trade-offs in counterfactual generation, distinguishing it from existing methods. To evaluate fairness, we introduce novel metrics for both group and subgroup level analysis that explicitly account for these trade-offs. Experiments on benchmark datasets show that FACEGroup effectively generates feasible group counterfactuals while accounting for trade-offs, and that our metrics capture and quantify fairness disparities.

**Keywords:** explanations · fairness · XAI · counterfactuals.

## 1 Introduction

AI-driven technologies increasingly shape critical decisions, making it essential to understand their underlying reasoning and evaluate their fairness. A variety of explanation methods have been proposed to enhance transparency [9,1], with counterfactual explanations (CFs) gaining prominence [34]. Individual CFs reveal how modifying specific features can alter model decisions, offering actionable insights. For example, consider a person whose loan application is rejected by a machine learning model; a CF might indicate that increasing annual income or reducing the debt-to-income ratio would lead to approval.

Prior work has primarily focused on individual counterfactual explanations (CFs) [18,16,10,33,27,29,25,31,4,2], with comparatively few studies addressing counterfactuals for groups of instances [28,23,19,20]. Group counterfactual explanations (GCFs) identify how a group of instances, often defined by shared characteristics or *protected attributes* such as sex or race, could collectively alter their features to achieve favorable outcomes. GCFs are not simply aggregations

of individual CFs; rather, they reveal common patterns or barriers affecting the group as a whole, which is critical for understanding systemic disparities and informing policy or organizational decisions. Previous studies introduce group-based approaches, by identifying common patterns among individuals with favorable outcomes [28], learning global translation vectors, and scaling them for GCFs [23], or constructing decision trees via stochastic local search [19]. In contrast, our work is the first to generate GCFs using a graph-based approach that enforces feasibility, supports subgroup-level analysis, and explicitly addresses the key trade-offs involved in counterfactual generation.

FACEGroup, our approach for generating Feasible and Actionable Group Counterfactual Explanations (GCFs), generates GCFs using a density-weighted feasibility graph [27], where nodes represent data points and edges denote feasible transitions that comply with real-world constraints. To ensure plausibility, we restrict connections to allow only small feature changes between data points. A key property of this graph is that feasibility constraints, cost limitations, and density weighting naturally partition the data into weakly connected components (WCCs), effectively dividing each group into subgroups with similar feasible counterfactual explanations.

The generation of group counterfactual explanations (GCFs) inherently involves balancing several key trade-offs: the proportion of factual instances within a group that are explained by the selected set of counterfactuals (coverage), the effort or change required for group members to achieve a counterfactual (cost), and the number of unique counterfactuals generated for the group (interpretability). To address these trade-offs, we introduce two algorithmic formulations based on the feasibility graph: the cost-constrained approach, which maximizes group coverage under a cost limitation, and the coverage-constrained approach, which minimizes the maximum cost required to achieve a specified coverage level. Both formulations are supported by mixed-integer programming solutions and greedy heuristics that operate at both the group and subgroup levels. Our approach also ensures that the generated counterfactuals remain feasible and actionable.

Finally, we introduce novel fairness metrics for group counterfactuals, which enhance existing fairness measures by capturing the various trade-offs in counterfactual generation and can be applied at both group and subgroup levels. We evaluate FACEGroup on real-world datasets, showing its effectiveness in fairness auditing. Compared to existing methods, FACEGroup produces more feasible and compact counterfactuals that align with the data distribution.

The rest of this paper is structured as follows: Section 2 formalizes the problem, Section 3 presents our algorithms, Section 4 introduces our fairness measures, Section 5 details experiments, Section 6 discusses related work, and Section 7 concludes.

## 2 Problem Definition

Let  $f : \mathbb{R}^d \rightarrow \{0, 1\}$  be a binary classifier which maps instances in a  $d$ -dimensional feature space into two classes, labeled 0 and 1. Let  $U \subseteq \mathbb{R}^d$  denote the in-

put space. A model prediction on an individual instance  $\mathbf{x} \in U$ , called *factual*, is explained by crafting a counterfactual (CF) instance  $\mathbf{x}' \in \mathbb{R}^d$  that is similar to  $\mathbf{x}$  but leads to a different outcome, i.e.,  $f(\mathbf{x}') \neq f(\mathbf{x})$  [34]. The changes in feature values from  $\mathbf{x}$  to  $\mathbf{x}'$  should be feasible and comply with real-world constraints, for instance, changes to immutable features, such as race or height, should be prohibited. Formally, a counterfactual  $\mathbf{x}'$  for  $\mathbf{x}$  is defined as:  $\mathbf{x}' = \arg \min_{\mathbf{x}'' \in \mathcal{A}_{\mathbf{x}}} \text{cost}(\mathbf{x}, \mathbf{x}'')$  s.t.  $f(\mathbf{x}'') \neq f(\mathbf{x})$ , where  $\text{cost}(\mathbf{x}, \mathbf{x}'')$  is a function measuring the cost of transitioning from  $\mathbf{x}$  to  $\mathbf{x}''$ . The *feasibility set*  $\mathcal{A}_{\mathbf{x}}$  denotes the set of counterfactuals attainable from  $\mathbf{x}$  via feasible changes.

It would be hard to trust a CF if it resulted in a combination of features that were unlike any observations the classifier has encountered before [34]. Therefore, CFs should also be coherent with the underlying data distribution. To ensure both feasibility and plausibility, we adopt a graph-based approach. Following [27], we construct a weighted directed graph  $G_U = (V, E, W)$ . Nodes correspond to instances in  $U$ , and an edge from node  $\mathbf{x}_i$  to node  $\mathbf{x}_j$  represents a feasible transition in the feature space. We call this graph *feasibility graph*. Transitions are further constrained by a cost threshold  $\epsilon$ , ensuring that only small-cost feature changes are allowed. This ensures that changes between instances are both feasible and small. The weight function  $W$  is defined using a density-based approach [27] to ensure that CFs lie in dense areas of the input space and avoid outliers. Each edge in  $G_U$  is assigned a weight  $W_{ij}$ , calculated as the product of the density of the instances around the midpoint of  $\mathbf{x}_i$  and  $\mathbf{x}_j$  estimated using a Kernel Density Estimator (KDE) [8], and the cost between instances:  $W_{ij} = KDE\left(\frac{\mathbf{x}_i + \mathbf{x}_j}{2}\right) \text{cost}(\mathbf{x}_i, \mathbf{x}_j)$ .

Given  $G_U$ , we now formally define the feasibility set  $\mathcal{A}_{\mathbf{x}}$  of factual  $\mathbf{x}$  as the set of instances  $\mathbf{x}'$  for which there is a path in  $G_U$  from  $\mathbf{x}$  to  $\mathbf{x}'$ , i.e., the set of instances that are reachable from  $\mathbf{x}$ :  $\mathcal{A}_{\mathbf{x}} = \{\mathbf{x}' \in \mathbf{U} \mid \mathbf{x}' \text{ is reachable from } \mathbf{x} \text{ in } G_U\}$ . These instances are the *feasible* CFs for  $\mathbf{x}$ .

Instead of finding a CF for a single factual  $\mathbf{x}$ , we are interested in providing CFs for a set  $X \subseteq U$  of instances mapped to the same class. Let  $X' \subseteq U$  be the set of instances mapped to the opposite class. Our goal is to identify a small subset  $S$  of  $X'$  of size  $k$  that best explains  $X$ . We limit the number of CFs to  $k$  for interpretability. To select  $S$ , we consider coverage-cost trade-offs. For a set of CFs  $S \subseteq X'$ , coverage is:

$$\text{coverage}(X, S) = |\{\mathbf{x} \mid \mathbf{x} \in X \text{ and } \exists \mathbf{x}' \in S \cap \mathcal{A}_{\mathbf{x}}\}|.$$

We overload the notation for *cost* to define the cost between an instance and a set, as well as between two sets:

$$\text{cost}(\mathbf{x}, S) = \min_{\mathbf{x}' \in S} \text{cost}(\mathbf{x}, \mathbf{x}'), \quad \text{cost}(X, S) = \min_{\mathbf{x} \in X} \max_{\mathbf{x}' \in S} \text{cost}(\mathbf{x}, \mathbf{x}').$$

The function  $\text{cost}(\mathbf{x}, \mathbf{x}')$  captures the cost of transforming  $\mathbf{x}$  to  $\mathbf{x}'$ , offering flexibility to adapt to specific problem requirements. For example, cost can be defined as the vector distance (e.g., L2 norm), the sum of edge weights along the shortest path in  $G_U$ , or simply the number of hops on this path. By emphasizing proximity in feature space and by considering dense paths, these definitions ensure

that the CFs are closely aligned with the data distribution. Our approach works with any definition of cost.

A necessary condition for  $\mathbf{x}'$  to be a feasible counterfactual for  $\mathbf{x}$  is that both  $\mathbf{x}$  and  $\mathbf{x}'$  belong to the same weakly connected component (WCC) of  $G_U$ . As a result,  $G_U$  induces a partition of the set of factual instances  $X$  into  $m$  disjoint subsets  $X_1, \dots, X_m$ ,  $m > 0$ . Each subset  $X_i$  contains instances in  $X$  that belong to the same WCC of  $G_U$  and thus share a common space of feasible counterfactuals, denoted  $X'_i$ , which also reside within the same component. This partitioning of  $X$  into subgroups with distinct feasible counterfactual spaces offers a meaningful perspective for analyzing model behavior at both the group and subgroup level, highlighting regions of the input space that support similar feasible explanations.

We now provide two definitions of the FACEGroup problem. Our first definition prioritizes cost over coverage, setting a threshold on cost, and our second definition prioritizes coverage over cost, asking for a set that provides a specified coverage degree  $c$ .

*Problem 1 (Cost-Constrained).* Given  $X, X', k \in \mathbb{N}^*$ , and cost threshold  $d \in \mathbb{R}_*^+$ , find  $S \subseteq X'$  with  $|S| \leq k$  and  $Q \subseteq X$  such that for every instance  $\mathbf{x} \in Q$  there exist an instance  $\mathbf{x}' \in S$  such that  $\text{cost}(\mathbf{x}, \mathbf{x}') \leq d$  and  $|Q|$  is maximized.

*Problem 2 (Coverage-Constrained).* Given  $X, X', k \in \mathbb{N}^*$ , and coverage degree  $c$ ,  $0 < c \leq 1$ , find  $S \subseteq X'$  with  $|S| \leq k$  such that  $\text{coverage}(X, S) \geq c|X|$  and  $\text{cost}(X, S)$  is minimized.

### 3 Algorithms

Our approach to generating feasible CFs is based on the feasibility graph  $G_U$ . Both optimization problems are NP-hard. The cost-constrained problem can be formulated as an instance of the maximum coverage problem, while the coverage-constrained problem is similar to the classical  $k$ -center problem [32].

In the following, we present two versions for both problems: (a) a global version that generates CFs for the whole set  $X$  and (b) a local version that generates CFs per subgroup  $X_i$ . We also show how the local version can be used to generate CFs for the whole group  $X$ . A common step in both problems involves computing, for each factual  $\mathbf{x}$ , the candidate counterfactuals, i.e., the feasibility set  $\mathcal{A}_{\mathbf{x}}$  and computing costs. To this end, we use Breadth-First-Search for vector costs (e.g., L2 distance) and Dijkstra’s algorithm for shortest path costs, with complexities of  $O(|V| + |E|)$  and  $O(|V| \log |V| + |E|)$ , respectively.

#### 3.1 The Cost-Constrained FACEGroup Problem

We solve this problem using two approaches: (a) a Mixed-Integer Programming (MIP) that explicitly models constraints for each factual-counterfactual pair while optimizing coverage, and (b) a Greedy approach that iteratively selects CFs to maximize coverage.

For the MIP solution of the global version of the problem, we define two binary decision variables. Let  $r_{\mathbf{x}\mathbf{x}'}$  = 1 if  $\mathbf{x}'$  covers  $\mathbf{x}$ ; and  $r_{\mathbf{x}\mathbf{x}'} = 0$ , otherwise, and  $u_{\mathbf{x}'}$  = 1 if CF  $\mathbf{x}'$  covers any instance in  $X$ , and  $u_{\mathbf{x}'} = 0$  otherwise. The goal is to maximize the number of covered factual instances:

$$\begin{aligned} \max \sum_{\mathbf{x}' \in X'} \sum_{\mathbf{x} \in X} r_{\mathbf{x}\mathbf{x}'} \quad \text{s.t.} \quad & \sum_{\mathbf{x}' \in X'} u_{\mathbf{x}'} \leq k \quad (1) \\ \sum_{\mathbf{x}' \in X'} r_{\mathbf{x}\mathbf{x}'} \leq 1, \quad \forall \mathbf{x} \in X \quad (2) \quad & r_{\mathbf{x}\mathbf{x}'} \leq u_{\mathbf{x}'}, \quad \forall \mathbf{x}' \in X', \forall \mathbf{x} \in X \quad (3), \\ & u_{\mathbf{x}}, r_{\mathbf{x}\mathbf{x}'} \in \{0, 1\}, \quad \forall \mathbf{x} \in X, \mathbf{x}' \in X'. \quad (4) \end{aligned}$$

While constraint (1) limits the number of selected CFs to at most  $k$ , constraint (2) enforces that each factual instance  $\mathbf{x}$  is assigned to at most one CF  $\mathbf{x}'$ . Constraint (3) guarantees that if a CF  $\mathbf{x}'$  is assigned to cover a factual instance  $\mathbf{x}$  ( $r_{\mathbf{x}\mathbf{x}'} = 1$ ) then  $\mathbf{x}'$  must be selected  $u_{\mathbf{x}'} = 1$ , and constraint (4) defines the binary decision variables. This formulation has  $O(2^{|X'|})$  complexity.

For the global Greedy version of the problem, we iteratively select counterfactuals (CFs) to maximize coverage. Let  $S_t$  be the set of counterfactuals selected at iteration  $t$ . We start with an empty set  $S_0 = \emptyset$ . At each iteration  $t$ , the algorithm selects the CF  $\mathbf{x}' \in X'$  that

$$\mathbf{x}' = \arg \max_{\mathbf{x}'' \in X'} (\text{coverage}(X, S_{t-1}) + \text{coverage}(X, \{\mathbf{x}''\})), \quad (5)$$

updates  $S_t = S_{t-1} \cup \{\mathbf{x}'\}$ , and terminates when either  $|S_t| = k$  or all instances in  $X$  are covered.

The worst-case complexity of this algorithm is  $O(k|X|)$ . Given the submodular nature of coverage, where the marginal gain of adding a new CF to the set  $S$  decreases as  $S$  grows, it adheres to the properties of submodular maximization. Consequently, the attained coverage is no worse than  $(1 - \frac{1}{e})$  times the optimal maximum coverage [17].

The Greedy algorithm can also be used to provide a counterfactual explanation for a subgroup  $X_i$  by applying it only to the corresponding WCC. We can also utilize this local version to provide counterfactuals for the whole group  $X$  by applying the Greedy algorithm iteratively to all  $m$  WCC as follows. Initially, we apply a single step of the Greedy algorithm at each WCC. Then, we select the CF that provides the best coverage and apply an additional step of the algorithm to the WCC from which the CF was selected. We repeat this until the maximum number  $k$  of counterfactuals is reached or all factual instances are covered. It is easy to see that this local version provides the same result as the global one. The local Greedy selection has the same complexity as the global Greedy approach, as it follows a similar process while iterating over WCCs, either scanning all  $|X'|$  candidates or evaluating coverage within each component.

### 3.2 The Coverage-Constrained FACEGroup Problem

To solve this problem, we employ two algorithms: a mixed-integer programming (MIP) and a Greedy 2-approximation algorithm [13]. While the Greedy

algorithm provides an efficient yet approximate solution, the MIP guarantees optimal results [7], but can become computationally expensive for large graphs.

For the MIP formulation, the solution is similar to the Cost-Constrained problem with the following modifications. The objective function minimizes the maximum cost  $d$  of the farthest instance while ensuring that  $\text{coverage}(X, S) \geq c|X|$ . Constraints (1), (2), (3), and (4) still apply, along with:

$$\sum_{\mathbf{x}' \in X'} \text{cost}(\mathbf{x}, \mathbf{x}') r_{\mathbf{x}\mathbf{x}'} \leq d, \quad \forall \mathbf{x} \in X \quad (6), \quad \sum_{\mathbf{x}' \in X'} \sum_{\mathbf{x} \in X} r_{\mathbf{x}\mathbf{x}'} \geq c|X| \quad (7).$$

Constraint (6) ensures that the cost of any node to its assigned center does not exceed  $d$ , enforcing the objective function, and Constraint (7) enforces that the desired coverage percentage is achieved. For full coverage,  $c = 1$ , constraint (2) becomes an equality constraint, and constraint (7) is no longer needed.

For the Greedy algorithm, the process begins by arbitrarily selecting the first counterfactual  $\mathbf{x}'$  and assigning all factuals  $\mathbf{x}$  within a cost of  $r$  to it, where  $r$  is initially set to the maximum cost between any factual and candidate counterfactual. We then iteratively select the counterfactual that is farthest from those already chosen and assign all factuals within a cost of  $r$  to it. This process continues until we reach the predefined coverage or the number of counterfactuals  $k$ . To find the smallest value of  $r$  that satisfies the coverage requirement, we employ a binary search. The complexity of this algorithm is  $O(k|X|\log(d))$ , since it assigns up to  $|X|$  factuals for each of the  $k$  selected counterfactuals and binary search adds this logarithmic factor  $\log(d)$ , where  $d$  is the range of costs considered.

Both the MIP and the Greedy approaches can be applied globally and locally. In the global version, we apply the algorithms on the  $G_U$  graph. In the local version, for a specific subgroup  $X_i$  of  $X$ , the algorithms are applied within the corresponding  $WCC$  of  $G_U$ .

We now describe how the local version can be used to solve the global version. Consider the case of full coverage ( $c = 1$ ) with  $m$  WCCs ordered arbitrarily as  $C_1, C_2, \dots, C_m$ . Achieving full coverage reduces to distributing  $k$  counterfactuals among these components. Since at least one counterfactual is required per WCC, the maximum allocation per WCC is at most  $k - m$ . First, we run MIP or Greedy within each WCC, varying  $k$  from 1 to  $k - m$ . Let  $l_i$  be the minimum counterfactuals needed to fully cover  $C_i$ . We start by assigning  $l_i$  to each  $C_i$ , then iteratively allocate remaining counterfactuals to the WCC with the highest cost until the total reaches  $k$ .

When  $c < 1$ , the task becomes more complex as we have to allocate both  $k$  and coverage  $c$  across the WCCs. Let  $F(1\dots i, k, n)$  be the minimum cost of allocating  $k$  counterfactuals that cover a total of  $n$  factuals considering connected components  $WCC_1, \dots, WCC_i$ , where  $n = c|X|$ . Similarly, let  $F(i, k, n)$  represent the minimum cost of allocating  $k$  counterfactuals to cover  $n$  factuals within component  $WCC_i$ . Then, we can solve the problem with time complexity

of  $O(m(kn)^2)$ , using dynamic programming as follows:

$$F(1\dots i, k, n) = \min_{1 \leq n' \leq n, 1 \leq k' \leq k} \{F(1\dots i - 1, k - k', n - n') + F(i, k', n')\}$$

For large graphs, solving the MIP at a global level can become computationally demanding, as the number of decision variables and constraints grows exponentially with the dataset size. To improve performance, we add constraints only for instances  $\mathbf{x}$  and  $\mathbf{x}'$ , such that  $\mathbf{x}' \in A_X$ , reducing unnecessary computations. For full coverage, the complexity of the global Greedy approach is  $O(|X|k \log(d))$  while the complexity for the local approach is  $O(m(k - m)|X_i|k \log(d))$ .

## 4 FACEGroup for Auditing Fairness

In this section, we examine algorithmic fairness through the lens of FACEGroup. Group fairness refers to a set of principles designed to ensure that protected groups, often defined by sensitive attributes such as gender, race, or age, are treated similarly by a classifier. Broadly, group fairness can be categorized into *demographic parity*, which requires that the proportion of positive outcomes reflects representation of the group in the population, and *error-based fairness*, which focuses on equalizing classification errors, such as false negative rates, across groups [36,9].

To audit fairness for a group  $X$ , we generate group counterfactual explanations (GCFs) for relevant subsets of  $X$ . For example, we generate GCFs for the negatively classified instances of  $X$  when auditing for demographic parity, or the false negatives of  $X$  when auditing for error-based fairness. Disparities in the GCFs generated for different groups (e.g., males vs. females) can reveal potential biases in the model.

Unlike existing approaches, FACEGroup supports *multi-level* fairness auditing by partitioning each group into subgroups according to the connected components of the feasibility graph. This allows us to examine unfair behavior not only at the group level, but also at the level of subgroups, offering finer-grained insight into patterns of bias. Furthermore, to capture the *key trade-offs* in generating counterfactuals, FACEGroup provides novel fairness metrics that are parameterized by the number  $k$  of counterfactuals, the cost  $d$ , and the coverage  $c$ . Introducing the number  $k$  in the fairness metrics allows for assessing interpretability, as groups requiring fewer CFs are more interpretable, it promotes trust, as models that require fewer CFs are more transparent, and it serves in detecting disparities in CF requirements across (sub)groups, factors previously overlooked.

**Burden-based Fairness Measures.** Counterfactuals provide a novel approach to measuring unfairness by evaluating both the disparities in outcomes between groups and the effort required by these groups to achieve fairness, i.e., to obtain the positive outcome. This effort, also called *burden*, is often estimated as the aggregated cost between the factuais in a group and their counterfactuals

[31,22]. However, measuring burden solely at the group level may obscure disparities within subgroups, as different subpopulations may face varying degrees of difficulty in achieving favorable outcomes.

We first define the minimum  $k$  ( $k_0$ ) and cost ( $d_0$ ) required for full coverage ( $c = 1$ ):

$$\begin{aligned} k_0 &= \min\{k \mid \exists S, |S| \leq k, \text{coverage}(X, S) = |X|\}, \\ d_0 &= \min\{d \mid \exists S, \text{cost}(X, S) \leq d, \text{coverage}(X, S) = |X|\}. \end{aligned}$$

Note that  $k_0$  is lower-bounded by the number of weakly connected components ( $k_0 \geq m$ ), and  $d_0$  does not exceed the largest WCC diameter.

We now introduce *AUC-based fairness measures* that assess trade-offs between cost, number of counterfactuals, and coverage of (sub)groups across a range of parameter values rather than at fixed points, avoiding biases from rigid parameter settings. The corresponding *saturation points* identify optimal thresholds for cost, number of counterfactuals, and coverage.

We define the set of counterfactuals  $S_{k,d}$  that maximize coverage under a cost constraint  $d$  as:

$$S_{k,d} = \operatorname{argmax}_{|S| \leq k, \text{cost}(X,S) \leq d} |\text{coverage}(X, S)|$$

and  $kAUC(k)$  as:

$$kAUC(k) = \int_{d_{min}}^{d_{max}} \text{coverage}(X, S_{k,d}) dd$$

that measures how efficiently a group can achieve coverage across a range of cost values for a given number of counterfactuals.

Similarly, we define  $dAUC(d)$  to evaluate how coverage improves as the number of counterfactuals increases under a fixed cost constraint, and  $cAUC(c)$  to quantify the effort required to reach a given coverage level by measuring the total cost over a range of counterfactual numbers. Figure 1 provides a visual representation of the AUC-based metrics.

There is also a minimum cost that provides the highest attainable coverage for  $k$ , we call it *saturation point* for  $k$  and denote it as  $sp(k)$ . Formally, it holds, for any  $d \geq sp(k)$ ,  $\text{coverage}(X, S_{k,d}) = \text{coverage}(X, S_{k,sp(k)})$ . Similarly, we define,  $sp(d)$  to determine the least number of counterfactuals needed to reach maximum coverage within a given cost constraint, and  $sp(c)$  to represent the minimum cost needed to achieve a desired coverage level, helping quantify the burden on different groups. Saturation points are shown in Figure 1.

**Attribution Measures.** FACEGroup also provides insights into feature importance by measuring how often a feature change is required to alter an outcome. Concretely, the *attribute change frequency (ACF)* metric captures how frequently a feature  $A$  changes between a factual instance  $\mathbf{x} \in X$  and its corresponding counterfactual  $\mathbf{x}' \in S$ :

$$ACF(X, S, A) = \frac{1}{|X|} \sum_{\mathbf{x} \in X} (1 - \delta(\mathbf{x}_A, \mathbf{x}'_A)),$$

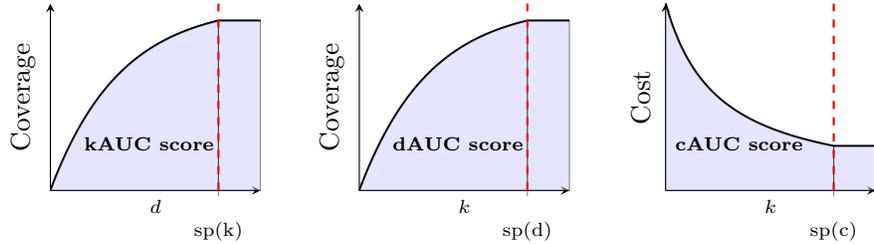


Fig. 1: AUC scores and saturation points

where  $\delta(\mathbf{x}_A, \mathbf{x}'_A)$  is the Kronecker delta, returning 1 if the feature remains unchanged and 0 otherwise. and  $\mathbf{x}_A$  and  $\mathbf{x}'_A$  represent the values of  $A$  in the factual and counterfactual instances, respectively. For each factual instance, we get the corresponding counterfactual instance with the minimum cost, i.e.,  $\mathbf{x}' = \operatorname{argmin}_{\mathbf{x}'' \in S} \operatorname{cost}(\mathbf{x}, \mathbf{x}'')$ .

## 5 Experimental Evaluation

The goal of our experimental evaluation is twofold: (a) to demonstrate the effectiveness of FACEGroup in fairness auditing and (b) to compare FACEGroup with baseline group counterfactual methods.

For fairness auditing, we use the widely studied `Adult`<sup>4</sup> dataset for income classification. To benchmark FACEGroup with baselines, we extend evaluations to additional datasets derived from US Census surveys, `AdultCA`<sup>5</sup>, `AdultLA`<sup>2</sup>, and other domains including `COMPAS`<sup>6</sup>, `Student`<sup>7</sup>, `German Credit`<sup>8</sup>, and `HELOC`<sup>9</sup>. Further details on preprocessing, parameter settings, and configurations, as well as additional experiments on other datasets, are in the supplementary material. The source code is available online<sup>10</sup>.

First, we construct the feasibility graph  $G_U$ . An edge exists from a  $\mathbf{x}_i$  to a  $\mathbf{x}_j$  if the transition from  $\mathbf{x}_i$  to  $\mathbf{x}_j$  is feasible and within threshold  $\epsilon$ . We use a small set of generic feasibility constraints prohibiting unrealistic modifications, such as changing the values of immutable attributes (e.g., race) or the directionality of others, such as decreasing the value of the age attribute. The full set of constraints used is in the supplementary material. We define groups based on the sensitive attribute *Gender*:  $G_0$  (females) and  $G_1$  (males).

Figure 2 depicts the impact of varying  $\epsilon$  on graph connectivity metrics, showing values up to the point where nearly all instances are connected, minimizing singleton nodes. Smaller  $\epsilon$  values result in sparser graphs, ensuring that connected instances are more similar, leading to more plausible, small-step transitions. Conversely, larger  $\epsilon$  values create denser graphs by incorporating connections between more distant instances, allowing for larger transition steps. To

<sup>4</sup> Adult   <sup>5</sup> Adult-CA-LA Datasets   <sup>6</sup> COMPAS   <sup>7</sup> Student   <sup>8</sup> German Credit  
<sup>9</sup> HELOC   <sup>10</sup> Project Repository

balance plausibility with connectivity, we select the smallest possible  $\epsilon$  that maintains a highly connected graph while minimizing singleton nodes. For the `Adult` dataset, we set  $\epsilon = 0.4$ . Further results for the selection of  $\epsilon$  on the remaining datasets can be found in the supplementary material.

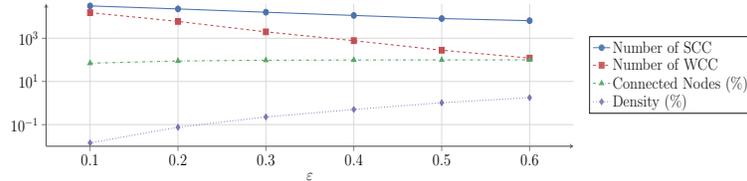


Fig. 2: Feasibility graph connectivity based on the  $\epsilon$  constraint.

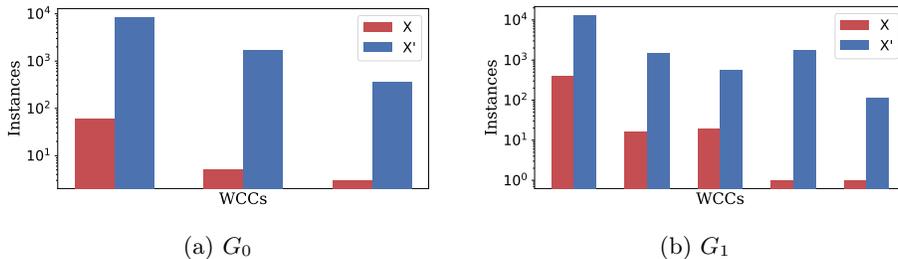
### 5.1 Auditing Fairness

In this set of experiments, we apply our algorithms to audit fairness. Without loss of generality, we focus on finding GCFs for the negatives for both groups  $G_0$  and  $G_1$ . We use an XGBoost classifier optimized via hyperparameter tuning. We consider only the instances in  $G_0$  and  $G_1$  for which at least one feasible candidate CF exists and use the  $L_2$  distance as the cost function.

**Burden Analysis.** A key strength of FACEGroup is its ability to uncover subgroup behaviors within the groups  $G_0$  and  $G_1$  through the feasibility graph  $G_U$ , which naturally partitions each group into WCCs, representing subpopulations that share feasible CF transformations. Figure 3 visualizes the distribution of factual instances ( $X$ , red) and feasible counterfactual candidates ( $X'$ , blue) across the subgroups (WCCs) of each group. We observe that  $G_1$  exhibits a more fragmented structure, with CFs more widely spread across subgroups compared to  $G_0$ , suggesting that  $G_1$  has a higher degree of variability in the transformations required for favorable outcomes. Table 1 depicts the minimum resources ( $k_0$  and  $d_0$ ) needed for full coverage per subgroup (WCC).  $G_1$  requires more CFs ( $k_0 = 12$ ) than  $G_0$  ( $k_0 = 9$ ) and higher minimum cost ( $d_0 = 1.04$ ) than  $G_0$  ( $d_0 = 0.93$ ), suggesting greater heterogeneity in the CF pathways needed for full coverage.

Analyzing subgroups is crucial, as group-level fairness assessments can mask heavily disadvantaged subpopulations, leading to misleading conclusions about the equitable distribution of the burden. At the subgroup level, the *Black* subgroups (that correspond to  $WCC_1$  in both groups) exhibit the highest  $k_0$  and  $d_0$ , indicating that they face greater barriers to obtain favorable decisions. Notably, the subgroups with the most factual instances also bear the highest burden, indicating a disproportionate impact on overall group difficulty.

Table 2 reports  $kAUC$ ,  $dAUC$ ,  $cAUC$ , saturation points  $sp$ , and the minimum, or maximum values for coverage and cost, that correspond to each  $sp$ .

Fig. 3: Distribution of  $\mathbf{X}$ ,  $\mathbf{X}'$  per  $WCC$  of the subgroups  $G_0$  and  $G_1$ .Table 1:  $k_0$  and  $d_0$  for each  $WCC$  of each group and overall for each group.

	WCCs										Overall	
	$WCC_1$		$WCC_2$		$WCC_3$		$WCC_4$		$WCC_5$		$k_0$	$d_0$
	$k_0$	$d_0$										
$G_0$	7	0.93	1	0.74	1	0.49	–	–	–	–	9	0.93
$G_1$	4	1.04	3	0.61	3	0.78	1	0.46	1	0.20	12	1.04

Scores are normalized by the optimal  $AUC$  per metric. Higher  $kAUC$ ,  $dAUC$  and lower  $cAUC$  are preferred.

For  $kAUC$ , saturation points ( $sp$ ) are expected to decrease as more CFs are provided. Initially, at  $k = 1$ ,  $G_1$  achieves higher maximum coverage, reflecting larger available transitioning costs, enabling more instances to be efficiently covered at low  $k$ . However, as the number of CFs increases,  $G_0$  reaches full coverage first, exhibiting better overall efficiency (higher  $kAUC$ ) and requiring fewer resources (lower  $sp$  values) compared to  $G_1$ . For  $dAUC$ , saturation points should decrease as higher-cost connections are allowed. At  $d = 0.1$ ,  $G_0$  has a lower  $sp(d)$ , indicating fewer feasible low-cost available transitions, compared to  $G_1$ . As cost increases,  $G_0$  effectively utilizes connections to reach full coverage with fewer CFs, while  $G_1$  requires higher costs to achieve maximum comparable coverage. However, when  $d \in [0.8, 1.5]$ ,  $G_1$  exhibits stronger coverage efficiency gains, suggesting  $G_0$  is more efficient at lower costs while  $G_1$  benefits more from cost relaxations. For  $cAUC$ , both groups experience similar cost burdens for achieving intermediate coverage levels 0.25, 0.5 and 0.75. However, at full coverage ( $c = 1.0$ ),  $G_1$  incurs significantly higher costs, as reflected in both  $cAUC$  and minimum cost. The consistently higher  $sp(c)$  values for  $G_1$  suggest that more CFs are required to reach cost-efficient solutions, reinforcing a systemic disadvantage in obtaining full coverage at minimal cost while maintaining interpretability.

**Attribution Analysis.** To further analyze subgroup disparities, we use the  $ACF$  metric per WCC, quantifying how often specific features are altered in CFs, providing insights into the different factors driving classification decisions. Figure 4 presents the frequency of modified attributes for each WCC of  $G_0$  and  $G_1$ ,

Table 2:  $kAUC$ ,  $dAUC$ ,  $cAUC$ , and saturation points.

Parameter	Value	$G_0$			$G_1$		
<b><math>kAUC</math> metrics</b>							
		$sp(k)$	Max Cov.	$kAUC$	$sp(k)$	Max Cov.	$kAUC$
$k$	1	1.1	63.08	0.50	1.3	65.75	0.54
	5	1.1	93.85	0.82	1.1	97.49	0.85
	9	1.1	100.0	0.90	1.1	99.09	0.89
	13	0.7	100.0	0.92	1.1	100.0	0.91
<b><math>dAUC</math> metrics</b>							
		$sp(d)$	Max Cov.	$dAUC$	$sp(d)$	Max Cov.	$dAUC$
$d$	0.1	6	12.31	0.10	12	12.78	0.08
	0.8	10	100.0	0.89	12	99.31	0.93
	1.5	9	100.0	0.93	12	99.77	0.95
	2.2	9	100.0	0.93	12	99.77	0.95
<b><math>cAUC</math> metrics</b>							
		$sp(c)$	Min Cost	$cAUC$	$sp(c)$	Min Cost	$cAUC$
$c$	0.25	12	0.14	0.10	20	0.12	0.11
	0.50	18	0.22	0.17	23	0.20	0.17
	0.75	22	0.28	0.25	25	0.30	0.25
	1.00	16	0.55	0.56	20	1.40	0.72

respectively, and shows that subgroup-specific variations exist in the importance of different features. For  $G_1$ , we include only the three largest WCCs, excluding those with few factual instances, as they lack representativeness. A common trend across all WCCs in both groups is that an increase in *age* is frequently required for a favorable outcome, suggesting that the model associates age with work experience or financial stability. Within  $G_0$ , the Asian-Pacific-Islander individuals ( $WCC_3$ ) require fewer modifications compared to the Blacks ( $WCC_1$ ) and Amer-Indian-Eskimos ( $WCC_2$ ) and do not rely on *relationship status* or *marital status*, unlike the others. In  $G_1$ , despite similar CF difficulty (Table 1), financial interventions differ: Amer-Indian-Eskimos ( $WCC_2$ ) require career-related changes (*employment status*, *occupation*, *education*), while Asian-Pacific-Islanders ( $WCC_3$ ) depend on increasing *capital gain*. More broadly, *capital gain* is largely absent from both groups of CFs except for  $G_1 - WCC_3$ , highlighting subgroup differences in financials to favorable outcomes. Finally, CFs in  $G_1$  rarely modify *relationship status*, unlike in  $G_0$ , where it is frequently altered. Instead, *educational* and *occupational* factors are highly important.

## 5.2 Comparison with Baselines

We evaluate FACEGroup against existing CF generation methods, specifically: (a) with FACE [27], a graph-based method for individual CFs, and (b) with AREs [28] and GLOBE-CE [23], two state-of-the-art GCF approaches.

**Comparison with Individual CFs** Given a group  $X$ , FACEGroup generates a small set  $S$  of  $k$  counterfactuals to cover  $X$ . To evaluate the efficiency

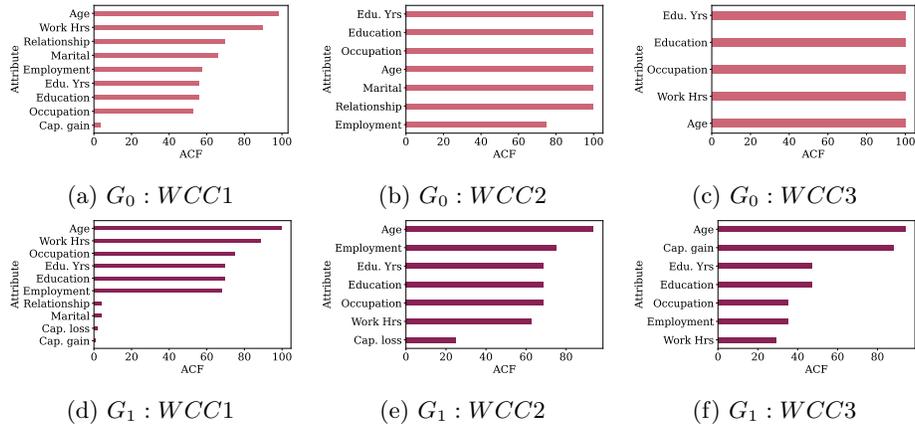


Fig. 4: ACF across the subgroups of each group

of this approach, we compare the associated cost with the cost of generating *individual counterfactuals* for each instance in  $X$ , which serves as a lower bound on the cost when the constraint on  $k$  is relaxed. For generating individual counterfactuals, we use FACE, since it is also based on a feasibility graph. For these experiments, we generate CFs for the full population  $G = G_0 \cup G_1$ . We assess how closely GCFs from FACEGroup approximate the optimal costs of individual CFs from FACE. First, we apply FACEGroup to generate the set  $S$  of CFs by solving the coverage-constrained problem. Then, we apply FACE to all factuals covered by  $S$  using the same cost function. As a cost function, we use both: (a) the weighted shortest path cost in  $G_U$  (originally used in FACE), and (b) the  $L_2$  distance.

Figure 5 shows the cost comparison for  $k$  CFs from 1 to  $k_0$  in 10 equal steps, with normalized costs. As expected, FACE achieves the lowest costs, while FACEGroup, which prioritizes group-level explanations, incurs slightly higher but still near-optimal costs. FACEGroup maintains near-optimal shortest path costs in datasets like **German Credit** and **HELOC**, where feasible transformations remain efficient. However, in **Adult**, costs increase due to the challenge of balancing feasibility with compact group CFs. Similar trends hold across other datasets, with full results and parameter details provided in the supplementary material.

**Comparison with GCF Methods** We compare FACEGroup with two state-of-the-art GCF baselines: AReS [28] and GLOBE-CE [23]. AReS mines frequent itemsets from individuals who achieved the desired outcome, selecting a small, interpretable set of rules via a submodular objective. GLOBE-CE defines global CFs as translation vectors applied to groups, scaling them across a range of values to adapt to individuals.

Both baselines without feasibility and plausibility constraints achieve at least 70% coverage. AReS generates 3 to 20 rules, while GLOBE-CE produces a significantly larger set, ranging from 10 to 612 CFs, due to the multiple scales on

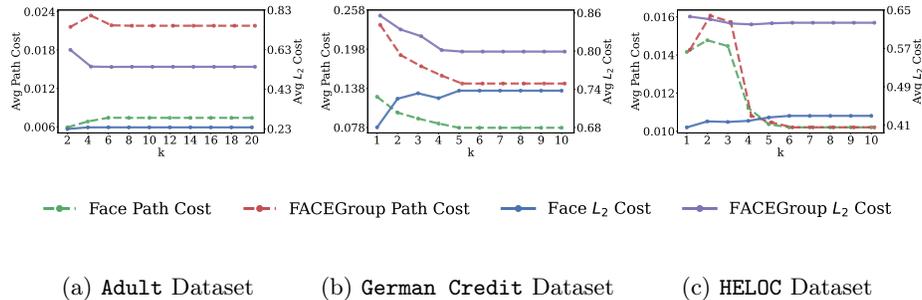


Fig. 5: Comparison of FACEGroup and FACE on average CF costs.

Table 3: Comparison with baselines.

Dataset	$\epsilon$	FC	AReS		GLOBE-CE		FACEGroup	
			$r$	Cov. (%)	$k$	Cov. (%)	$k$	Cov. (%)
Adult	0.4	all	18	15.68	421	0.24	21	100
	0.4	none	18	52.26	421	84.56	10	100
AdultCA	0.7	all	20	11.36	612	11	133	100
	0.7	none	20	11.36	612	11.50	15	100
AdultLA	0.5	all	20	12.9	342	12.9	59	100
	0.5	none	20	23.11	342	22.63	13	100
Student	3.0	all	3	33.3	10	50	3	100
	3.0	none	3	75	10	66.67	2	100
COMPAS	0.3	all	20	11.85	124	20	13	100
	0.3	none	20	16.3	124	25.93	13	100
German Credit	2.9	all	4	0	18	26.32	6	100
	2.9	none	4	42.11	18	73.68	2	100
HELOC	1.4	all	11	1.98	74	1	4	100
	1.4	none	11	71.29	74	72.28	2	100

top of the translation vectors. Detailed results are in the supplementary material. To assess feasibility, we integrate CFs into the feasibility graph  $G_U$  and measure feasibility coverage as the proportion of CFs with at least one feasible transition. We analyze this under all feasibility constraints and a relaxed setting with only the plausibility constraint  $\epsilon$ . Table 3 highlights the limitations of baselines: with full constraints, AReS and GLOBE-CE remain below 50% feasibility coverage, indicating that many CFs violate real-world constraints. In contrast, FACEGroup achieves 100% feasibility coverage with a compact CF set. Relaxing constraints improves coverage for baselines, particularly for GLOBE-CE, which benefits from its low-cost translation vectors. However, FACEGroup still maintains full feasibility coverage with fewer CFs, demonstrating its ability to generate feasible, actionable CFs without sacrificing interpretability or plausibility.

## 6 Related Work

Explanations have become central in machine learning research [9,14], particularly in high-stakes domains such as healthcare and education. Among various explanation methods, CFs have gained prominence for their ability to reveal actionable changes leading to a desired outcome. Wachter et al. [35] first formulated CFs as an optimization problem, minimizing the cost between an instance and its CF while ensuring a prediction change. Subsequent work [15,25,18,12,34,31,27] refined CF generation, emphasizing properties such as feasibility, actionability, sparsity [34], and robustness [16]. Several approaches optimize CF search using genetic algorithms [31,10], integer programming [30,33], and cost-based heuristics [12].

FACE [27] constructs a density-weighted feasibility graph where counterfactuals are generated via shortest paths in the graph, focusing on individual explanations that balance proximity and data manifold alignment. While FACEGroup builds on this graph structure, and further introduces three key innovations: (1) multi-level subgroup analysis, where WCCs of the feasibility graph naturally partition groups into interpretable subgroups with shared feasibility constraints, (2) GCF trade-off-aware algorithms, rather than relying on individual shortest-path searches, and (3) cost function agnosticism.

While most methods focus on individual CFs, recent work explores GCFs for multiple instances. AReS [28] defines subgroup-specific CF rules, optimizing for correctness, coverage, cost, and interpretability. GLOBE-CE [23] learns global translation vectors, applying them at different scales to generate CFs that maximize coverage. CET [19] uses decision trees for group actions to enhance transparency and consistency, while mixed-integer programming has been used to optimize collective CFs under linking constraints [5]. CounterFair [21] generates fair GCFs by selecting a subset via mixed-integer programming to balance cost and fairness. Unlike these approaches, FACEGroup enforces feasibility constraints, ensuring GCFs adhere to real-world constraints. Most group-based methods only prevent changes in sensitive attributes but lack directional constraints, leading to CFs that may violate plausible transformations. Notably, GLOBE-CE selects random feature perturbations, which can result in unrealistic CFs. In contrast to these methods, FACEGroup generates CFs at both group and subgroup levels, systematically handling the trade-offs in CF generation.

Explanations are utilized to assess algorithmic fairness [9], ensuring decisions are not influenced by protected attributes [26,24,11,6]. Several CF-based approaches have been proposed to quantify fairness by measuring the burden quantified as the difficulty individuals face in achieving a favorable outcome per group [22,31,12,20,28]. Methods like [31,22] generate individual CFs and calculate burden per group as the average sum of pairwise costs to assess fairness. PreCoF [12] distinguishes between explicit bias, when individual counterfactuals require changes only in sensitive attributes, and implicit bias, when, after removing sensitive attributes from model training, other features disproportionately influence different groups. [28,23] suggest that generated rules and global translation vectors can be used to manually audit for unfairness in subgroups of

interest. FACTS [20] builds on AReS and introduces burden-based fairness metrics, but evaluates fairness only under specific settings. For instance, its Equal Cost of Effectiveness metric compares the minimum cost needed for protected subgroups to reach a fixed aggregate effectiveness level, defined as the proportion of individuals able to achieve the desired outcome via counterfactuals. In contrast, our burden-based fairness metrics assess disparities across a range of costs, coverage levels, and numbers of counterfactuals, offering a more comprehensive perspective that captures potential disparities across various combinations of these factors. Unlike the other approaches, FACEGroup introduces fairness metrics that assess fairness at both group and subgroup levels, explicitly accounting for trade-offs between cost, coverage, interpretability, and feasibility.

## 7 Conclusions

In this paper, we propose FACEGroup, a novel graph-based framework for group counterfactual generation that addresses limitations in existing methods by incorporating real-world feasibility constraints and managing trade-offs in counterfactual generation. We also introduce novel fairness measures that allow auditing fairness both at the group and subgroup levels, offering insights on the trade-offs between cost, the number of generated counterfactuals, and coverage. In future work, we plan to extend the use of the feasibility graph to define path-based fairness metrics. We also aim to adapt our approach to multi-class classification and regression settings.

## 8 Acknowledgment

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